



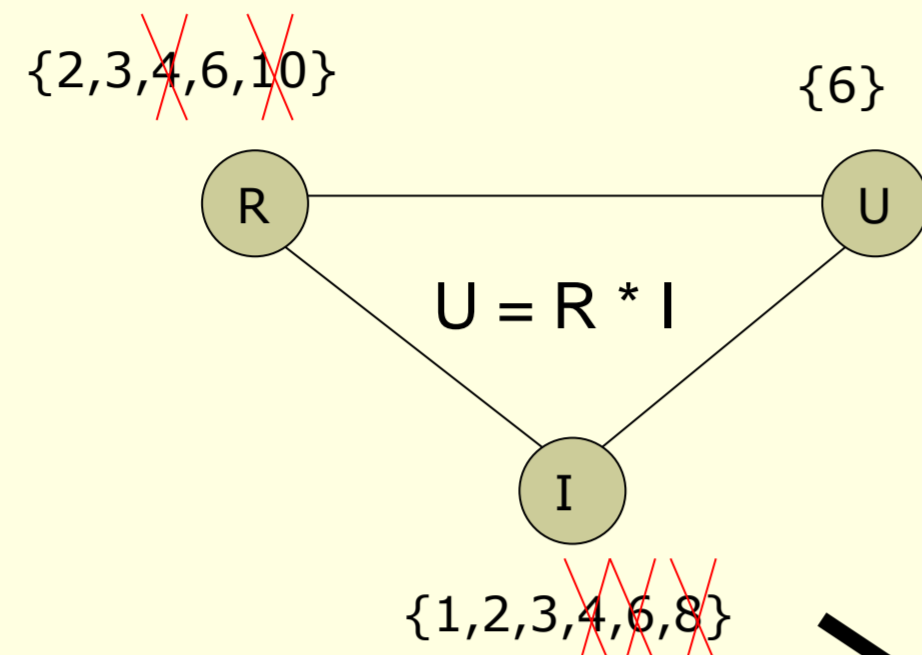
CAVERN:

Constraints and Abstractions for program VERification



Constraint Programming

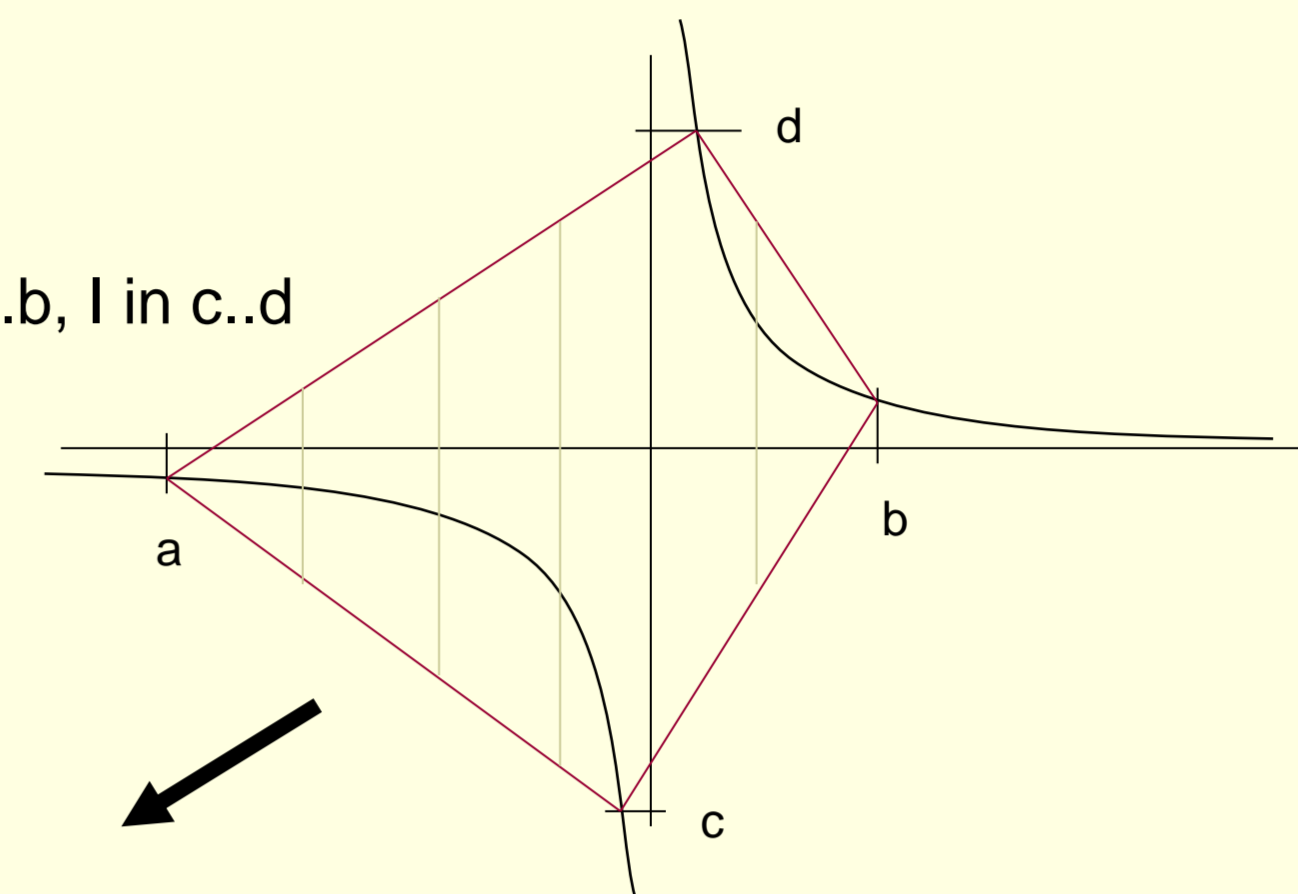
Exploit relations (constraints) to infer new informations on objects that represent unknowns (variables)



$$1 = R * I, \quad R \text{ in } a..b, I \text{ in } c..d$$

Abstractions

Over-approximate the computation of relations to benefit from powerful solving techniques (e.g. Linear Programming)



Program Verification

P1a	/*@ behavior P1a : assumes flag = 0 && n = 8 && forall i in 1..8, ar[i] = cos(2*M_PI*i/n); ensures \result == 0;
P1b	/*@ behavior P1b : assumes flag = 1 && n = 8 && forall i in 1..8, ar[i] = cos(2*M_PI*i/n); ensures \result == 0;
...	

Are properties P1a,P1b, ...
verified
by this implementation ?

```
int fft1(int n, int flag) {
  int i, j, k, it, xp, xp2, j1, j2, iter;
  double sign, w, wr, wi, dr1, dr2, di1, di2, tr, ti, arg;

  if(n < 2) return(999);  iter = log((double)n)/log(2.0);  j = 1;
  for(i = 0; i < iter; i++)  j *= 2;  if(fabs(n-j) > 1.0e-6) return(1);
  sign = ((flag == 1) ? 1.0 : -1.0);  xp2 = n;
  for(it = 0; it < iter; it++) {  xp = xp2;  xp2 /= 2;  w = PI / xp2;
    for(k = 0; k < xp2; k++) {  arg = k * w;  wr = cos(arg);  wi = sign * sin(arg);  i = k - xp;
      for(j = xp; j <= n; j += xp) {
        j1 = j + i;  j2 = j1 + xp2;  dr1 = ar[j1];  dr2 = ar[j2];  di1 = ai[j1];  di2 = ai[j2];
        tr = dr1 - dr2;  ti = di1 - di2;
        ar[j1] = dr1 + dr2;  ai[j1] = di1 + di2;  ar[j2] = tr * wr - ti * wi;  ai[j2] = ti * wr + tr * wi;
      } }
    j1 = n / 2;  j2 = n - 1;  j = 1;
    for(i = 1; i <= j2; i++) {
      if(i < j) { tr = ar[j-1];  ti = ai[j-1];  ar[j-1] = ar[i-1];  ai[j-1] = ai[i-1];  ar[i-1] = tr;  ai[i-1] = ti; }
      k = j1;  while(k < j) { j -= k;  k /= 2; }  j += k; }
    if(flag == 0) return(0);  w = n;
    for(i = 0; i < n; i++) { ar[i] /= w;  ai[i] /= w; }
    return(0); } }
```

WP2: Memory models

Fine-grained memory model
of C and Bytecode Java

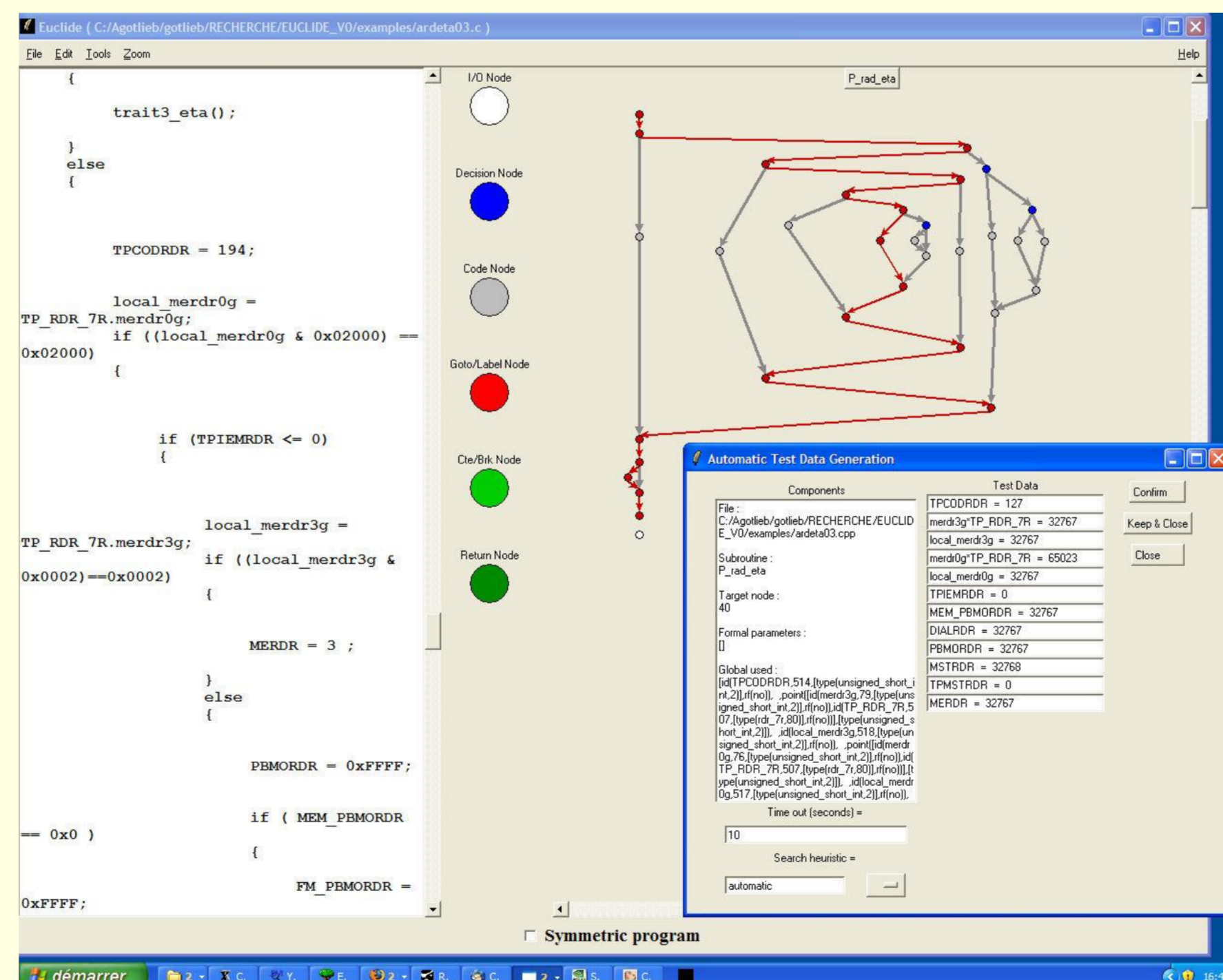
PhD of Florence Charreteur (Mar. 2010)
PhD of Mickael Delahaye (Oct. 2011)

F. Charreteur, B. Botella, and A. Gotlieb.
**Modelling dynamic memory management
in constraint-based testing.**
The Journal of Systems and Software,
82(11):1755–1766, 2009

F. Charreteur, A. Gotlieb
**Constraint-based test inputs generation for J
ava Bytecode.**
ISSRE 2010

M. Delahaye, B. Botella, A. Gotlieb
Explanation-based generalization of infeasible path.
ICST'10, Paris.

WP3: Integers and iterative computations



A. Gotlieb, B. Marre, and M. Leconte.
Constraint solving on modular integers ModRef'10

Olivier Ponsini, H el ene Collavizza, Carine F ed ele, Claude Michel, Michel Rueher,
Automatic Verification of Loop Invariants ICSM 2010

WP4: Floating-point computations

Properties over floating-point data
can be exploited to prune the search
Space, e.g.

Soit $z \in \mathbb{F}_{p,+\infty}$, avec $+\infty > z > 0$ et

$$z = 1.b_2 \dots b_i 0 \dots 0 * 2^{e_z} \text{ avec } b_i = 1$$

Soient

$$y = 1.1 \dots 1 * 2^{e_y + nb_z} \text{ avec } nb_z = p - i$$

$$x = y \oplus z$$

Alors

- $x \ominus y = x - y = z$,
- et il n'y a pas d' $x' \in \mathbb{F}_{p,+\infty}$, $x' > x$ ou d' $y' \in \mathbb{F}_{p,+\infty}$, $y' > y$ tel que $x' \ominus y' = z$

PhD of M. Said Belaid
Post-doc Matthieu Carlier

Bruno Marre, Claude Michel,
**Improving the floating point addition and
subtraction constraints** CP'10

Mohammed Said Belaid, Claude Michel, Michel Rueher,
**Approximating floating-point operations to verify
numerical programs** SCAN 2010,

Matthieu Carlier, A Gotlieb
Filtering by Maximum ULP ICTAI'2011

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